Introduction and Terms

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Semantic Technology for
Intelligence, Defense, and Security

October 26, 2010
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Topics

• Quick Intro: What is an ontology?
• Design of a toolset for realizing ontology
• Ontology development best practices
• Semantic federation best practices
• Reification for provenance and security
What is an ontology?
a logical domain model
reflects the real world
• Bullets

– More bullets
• Bullets

- More bullets
Bullets

- More bullets

machine readable
• Bullets

- More bullets

machine readable
- Bullets
  - More bullets

- Machine readable
machine readable
greater than the sum of its parts
greater than the sum of its parts
Abstract
Medical science often suffers from having so few data and so much imperfect knowledge that a rigorous probabilistic analysis ... is seldom possible. Physicians nevertheless seem to have developed an ill-defined mechanism for reaching decisions despite a lack of formal knowledge regarding the interrelationships of all the variables that they are considering. This report proposes a quantification scheme which attempts to model the inexact reasoning processes of medical experts...utilized in a rule-based computer diagnostic system. One such system, a clinical consultation program named , is described in the context of the proposed model of inexact reasoning.
A model of inexact reasoning in medicine

Edward H. Shortliffe
Bruce G. Buchanan

Mathematical Biosciences
Volume 23, Issues 3-4, April 1975, Pages 351-379

Abstract
Medical science often suffers from having so few data and so much imperfect knowledge that a rigorous probabilistic analysis ... is seldom possible. Physicians nevertheless seem to have developed an ill-defined mechanism for reaching decisions despite a lack of formal knowledge regarding the interrelationships of all the variables that they are considering. This report proposes a quantification scheme which attempts to model the inexact reasoning processes of medical experts...utilized in a rule-based computer diagnostic system. One such system, a clinical consultation program named MYCIN, is described in the context of the proposed model of inexact reasoning.
Why do we bother?

• We believe that data-intensive computer systems employing the discipline of ontology are more
  – Accurate
  – Flexible
  – Reusable
  – Responsive
  – Knowledge capture

• and our experience bears this out.
Sampling of recent projects

• Five Ongoing Intelligence Agency Projects
  – Analysis
  – Enterprise Ontology
  – Network Analysis
  – Network Health and Status
  – Determining ROI on Intelligence Collection

• University of New Haven/Institute for the Study of Violent Groups (ISVG)

• FreedTV
  – Social Networks-based recommender for entertainment content.

• Purus Technologies
  – semantic federation to support complex decision making.
How we build ontology

- Modified first-order logic
- Statements & entailments
- Knowledge Frame Language (KFL) & ECLIF
  - properties, n-ary relations, constants
  - function terms
  - rules and integrity constraints
Tools Overview

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October 26, 2010
Overview Goals

- High level product overview
- Examine capabilities of HighFleet Suite across several different user roles
- Point out usage patterns that will be most productive
What's a “HighFleet Suite”? 

- Ontology Development Support
  - KFL parsing, deploying; small data load, querying.

- Server Management
  - Start/Stop etc.

- Scalable extensional server
  - Billions of assertions
Architecture Overview

- Model Creation
- Data Store
- Inference Engine
- Scalability
- Interoperability
Model Creation

- KFL Ontology Modeling language, derived from the KIF first-order calculus effort.
- Used as syntactic sugar over Common Logic variant (ECLIF)
- Encompasses type and rule definition with extensions for measurement and time
- GUI deploy wizard and command line compilation support
Data Store

- Purpose-built hyper-graph store, leveraging an array of RDBMS technologies
- More a graph storage layer than triple store
- Optimized for concurrent, ACID access
- Allows 'what if' exploration through transactional visibility
- Range of scale-out options
Inference Engine

- Common Logic based FOL rules engine
- Static Semantic, In-Process Semantic, and Physical query optimizations
- Heavily optimized for parallelized execution
- Deep caching, multiple attack strategies
- Baked-in transactions, time handling, measurement
- Soup-to-nuts integration with the data store
Scalability – Rules

- Optimized for dense logic, not non-productive taxonomies and type hierarchies
- However, reasons efficiently over large taxonomies by treating terms as terms rather than properties.
- Rule-bases with thousands of predicates are fine
- There are some limits, but they tend to be sensible
- Distinction between schema and extensional data
Scalability – Extensional Data

- Distinguishing between schema and data, emphasizing scaling out extensional data
- Use RDBMS's for what they are good at: ACID and Indexing
- Ignore RDBMS modelling conventions
- Allow for fine-grained index control and statistics gathering
Scalability – Query Processing

- Vertical scalability with complete 64-bit support for large hardware situations
- Horizontal scalability by sharing an extensional store across multiple load-balancing query engines
- Intra-engine scaling with intelligent query program overlap reuse.
Runs (and tested!) across a variety of platforms, 64/32 bit, Windows, Linux, Mac OSX

Java API, JDBC-like.

REST-ful interface with complete access via HTTP/JSON protocol
The HighFleet Suite IODE

- Basic Administration
- Ontology Development
- Manipulating Data
- Advanced Administration
- Miscellaneous
Ontology Model Development

- KFL Model → Deployed XKS
- Use Non-Persistent XKS for Ontology Development
  - Fast for testing, low overhead
- Use Postgres/Oracle for serious deployments
Ontology Development

Browsing & Ontodoc

- Database → Browser
  - Traverse a loaded ontology
  - Show type hierarchies, instances, graphs of relations, etc.
  - Eminently searchable

- Ontology Doc Generation
  - Generates HTML description of a loaded ontology
Ontology Development

Types of XKS's

- Ontology Development XKS – Non Persistent
  - Designed to give exactly the same behavior as the more expensive standalone XKS servers
  - IODE license is all you need

- Persistent XKS
  - Licensed individually by assertion size
  - H2/Postgres/Oracle
Quote – Types of XKS's

• “The idea is that the non-persistent XKS and the standalone scalable XKS perform (in terms of correctness) exactly the same, since they share the same core engine with its highly parallel architecture.”
Extensional Data Overview

- Data loaded in terms of the model
- Scalable, performant, persistent (license dependent)
- Tools: Query tool, Questionator, Asserter
“Remember, creating a model in a vacuum is interesting, but any model is far more useful if you can express data in terms of that model in a scalable manner. The HighFleet Suite is designed to give you a model development environment that works as well for models with tiny data as large data.”
Querying

- Query Tool
  - Text query based

- Questionator
  - Graphical UI for building/saving queries
  - Graphing of results

- In both cases results can be exported
Fact Asserter

- Used to add data to an XKS
  - Add single facts interactively or files in bulk
- Generally, machine generated data might not need IC checking
- It is often useful to drop/restore indexes around a really large bulk load
Delete Assertions

- GUI Fact Deleter supports deleting
- Non-Persistent XKS is designed to mostly let you ignore deletions
- Usually, you don't delete, but instead denote with a time span
More *Interesting* Administration

- Backup/Export (data + schema or data only)
- Import Data, Restore Database
- Rename, Delete, Clone
- Configure
Exposing processing to the user: PseudoRel and PseudoAssertion

- Mostly for HighFleet internal use, they can be useful for debugging
- PseudoRel (showCompilationInfo (foo ?x)) will dump the query plan to the log as the query is run
- PseudoAssertion (owi.RelationIndex ?rel ?pos ?indexed) drops/creates indexes (useful for large ingest cases)
Future Features - Incremental

- Query optimization improvements
- More forgiving error reporting
- More work on web interface
- MORE
Future Features – Large Scale

- 64-bit off-heap caching, clustered caching
- User-controlled materialization
- Mutable (forward) ontology modifications
- eclipse/IDE based structured KFL editing
- Multi-lingual logic processing
Developing a User-facing Ontology Application, Soup to Nuts

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October 26, 2010
Starting Points

- Data
- Loose query requirements
- Interface plans
- Business goals
Demo: Keshvad application

• Data
  – Movements and Communications of Persons of Interest (who, where & when)
  – Simple geographic data
  – Parents, spouses, group membership

• Requirements
  – Social network analysis
  – Support sophisticated target tracking
  – Fully exploit the data
Best practice for ontology development: ontological layers

**Interface**
Describes the end use of the system (sample queries, display)

**Application-specific domain**
Describes the world from a particular (enterprise’s) perspective

**Application-independent domain**
Concepts useful to many perspectives

**System support**
Basic concepts like relation
Best practice: reuseable upper level ontology

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Highfleet Upper and Mid-Level Ontologies (ULO and MLO)
Best practice: reuseable upper level ontology

- Property, Material Role
- Relations, Subsumption
- Dates, numbers, strings
- Binary Relation, Transitive Closure BR
- Event
- gt, lt
Best practice: Leverage industry-standard domain models

- **Interface**: Describes the end use of the system (sample queries, display)

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Use existing standards and best practices (even if it hurts a little)
Best practice: Leverage industry-standard domain models

- Region Connection Calculus (RCC)

Model Development Task

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Keshvad ontology

**Interface**
target tracking & network visualization

**Application-specific domain**
movements & communications

**Application-independent domain**
social networking & geography (RCC)

**System support**
Deep dive: exploiting a movement event

1. There was a movement.
   (Movement trv-02)

2. This movement occurred from Jan 24 to Jan 26.
   (occursAt trv-02 (Span (Date "2008 Jan 24") (Date "2008 Jan 25")))

3. This movement was from the city Shiraz.
   (from trv-02 (keyCity "Shiraz" IR "Fars"))

4. This movement was to the city Zahedan.
   (to trv-02 (keyCity "Zahedan" IR "Sistan e Baluchestan"))

5. Keshvad made this movement.
   (agent trv-02 id-00)

6. The city Zahedan is located in the country Iran.
   (locatedIn (keyCity “Shiraz” “IR” “Fars”) (keyCountry “IR”))

What can you conclude from these facts?
Deep dive: exploiting a movement event

- Where was Keshvad located at the start of Jan 24th? At the end of Jan 26th?

(=> (and (Movement ?e) (from ?e ?place) (agent ?e ?a) (hasStartTime ?e ?t)) (holdsIn ?t (locatedIn ?a ?place)))

If you move from some place, then at the start of that movement you are located at that place.

- Similar rule for to
Deep dive: exploiting a movement event

- When was Keshvad in Iran?
- \( (\text{TransitiveBR locatedIn}) \)

\textit{If a is located in b, and b is located in c, then a is located in c.}
Deep dive: exploiting a movement event

• When was Keshvad in Iran?

• (TransitiveBR locatedIn)

  If $a$ is located in $b$, and $b$ is located in $c$, then $a$ is located in $c$.

  ($\Rightarrow$ (and (locatedIn ?a ?b)
   (locatedIn ?b ?c))
   (locatedIn ?a ?c))
Deep dive: exploiting a movement event

(schema (TransitiveBR ?Rel)

  (=> (and (?Rel ?a ?b)
           (?Rel ?b ?c))
       (?Rel ?a ?c)))

  (=> (and (holdsIn ?t (locatedIn ?a ?b))
           (holdsIn ?t (locatedIn ?b ?c)))
       (holdsIn ?t (locatedIn ?a ?c))))
Deep dive: exploiting a movement event

• It’s Jan 28th. Where was Keshvad last seen and how long ago?

(=>
  (and
    (Agent ?Agent)
    (Time ?QueryTime)
    (maxTermf ?T (?T ?Y) (and
      (atOrBefore ?T ?QueryTime)
      (observedIn ?T ?Agent ?Y))
      ?MaxTime)
    (observedIn ?MaxTime ?Agent ?Region)
    (ifThenElse
      (temporalIntersection ?MaxTime ?QueryTime ?Int)
      (= ?Elapsed (hours 0))
      (and
        (endingInstant ?MaxTime ?Left)
        (startingInstant ?QueryTime ?Right)
        (duration (Span ?Left ?Right) ?Elapsed)))
    (lastSeenIn ?QueryTime ?Agent ?Region ?Elapsed)))
Best practice: test-driven development

- Sample data set & test cases
- Ideally working w/ interface developer or from client req’s
- Store test cases as ontological data for easy handoff
Best practice: test-driven development

- Show test queries and run a few
Best Practices for Client Application Development
Semantically Integrating Legacy Databases

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October 26, 2010
What are we trying to do?

- We want to convert multiple structured data sources
  - different formats
  - internal and external
  - constantly changing

- into new, actionable information
  - could not be discovered from a single source
  - does not exist as the intersection of row and column
  - needs to be current

- within time and budget constraints.
  - a few weeks is not a reasonable query response time
  - minimize use of expensive resources
Best practice: create a logical model of each data source

Problem: Existing data federation processes are time-consuming and brittle.
Solution: Generate logical models of sources & map them to a single ontology.

- Central ontology means less mapping, **lower cost, faster implementation**.
  - “Every lab has a different LIMS, but they all use the same chemistry.”
  - Lower cost to add new data sources.

- Separates the logic of integration from the implementation of integration for **greater flexibility, easier maintenance**.
  - Traditional integration tools require you begin with the end in mind.
  - Semantic approach is more responsive to schema and requirement changes.
  - Allows you to build the model incrementally.
Best practice for ontology development: ontological layers

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**System support**
Basic concepts like relation
SemFed adds layers

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**Connection**
Describes how data interrelates

**Base**
Logical model of each data source

**System support**
Basic concepts like relation
Automatic Data Source Analysis

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- **Base**
  Logical model of data sources

- **System support**
  Basic concepts like relation

Generated by SemFed Toolkit
Metadata analysis and model generation

• If data is messy, clean it up
  – Primary key detection
  – Type inference

• If data is orderly, dig out every nugget of info
  – Aspect relations
  – Modeling recommendations

• Any structured data source

• Configure and iterate
SemFed Toolkit: KFL Generator

- Run KFL Generator against the CIA World Factbook
Metadata analysis and model generation

- View some generated KFL
- Mapping varies with source type
- Contains sufficient information to reverse engineer the source!
Demo: Keshvad ontology with remote data

- Collected open-source data about activities of violent groups (who, where & when)
  - 937 tables

- CIA World Factbook
  - 33 tables or as RDF
Best practice: Apply “thin coats” of connection ontology

- Mapping properties, mapping relations
- Extra rigor now = extra flexibility later
- Per-source mapping

- View some connecting ontology
- Domain ontology extensions were minimal
SemFed adds layers

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Logical model of each data source

**System support**
Basic concepts like relation
Problem

• Source 1
  – (keyCityTBL 2821440)
  – (locatedIn (keyCityTBL 2821440) Incident-01)

• Source 2
  – (keyCity Tehran IR)
  – (locatedIn Keshvad (keyCity Tehran IR))
Best practice: Use ontology to find equivalent entities

- Establish equivalence sets
- Materialize equivalence sets
- Enforce equivalence sets at the base level
Optimization for SemFed

- Ingest/extract static/slow-changing data
- Adjust caching time-to-live
Maintenance

• Accommodating a schema change
• Accommodating a requirements change
Reification and proofs in ECLIF

Support of advanced metadata applications in the XKS

Bill Andersen, Highfleet, Inc.

STIDS 2010 Tutorial Session
October 26, 2010
Customers ask about...

- Multi-level security filtering
- Discretionary access control
- Provenance
- Modularity
- Contextual reasoning
- ...
- We’re trying to do something about that!
Goals of this presentation

• Present our plans for support of advanced metadata applications in ECLIF / XKS
• Compare with RDF/OWL reification
• Explain our approach
• Solicit your input
  • What do(n’t) you like?
  • What haven’t we thought of?
  • Would you use this? For what applications?
Intended audience

( in order )

• Software engineers who want implement metadata in ontology-based applications

• Ontologists and technologists who are interested in applications of meta-level reasoning

• Project leads who need to understand the advantages of using meta-data in knowledge intensive applications
Outline

• Metadata and its applications
  • Reification and metadata
• Reification in RDF
• Reification in ECLIF
• Extension of reification to inference
• Examples and applications
Metadata

meta-level data

about your data

object-/domain-level data

about your domain
Why we need metadata

- Tracking history of data
- Transaction time, provenance
- Controlling data access
- Multilevel security, privacy
- And more ...

- To do this we have to talk about data
Reification

• ... is a mechanism for making sentences into objects of reference – talking about data

source([likes(bob, pizza)], joe)

• ... is the way we refer to data so we can use metadata to talk about it

• ... has lots of ways to approach it with various ramifications
Reification in RDF
Reification in RDF

an RDF triple $T$:  

$$<\text{bob \ likes \ pizza}>$$

and its reified form $T_R$:

$$\{ <s1 \ rdf:type \ rdf:Statement>,$$
$$<s1 \ rdf:subject \ \text{bob}>,$$
$$<s1 \ rdf:predicate \ \text{likes}>,$$
$$<s1 \ rdf:object \ \text{pizza}> \}$$

The reified form, $T_R$, of a triple $T$ is a set of four triples made by introducing a new name for $T$ and relating it to $T$'s constituent parts by special RDF vocabulary.

Usually, blank names are used for the new name.
Note on RDF semantics

http://www.w3.org/TR/rdf-mt/

• RDF semantics are based on graphs
• An RDF graph is a set of triples
  • This is important as we’ll see later
• RDF entailments are relativized to graphs
  • No inter-graph entailments
• Same goes for OWL / RDFS
RDF reification semantics - I

- RDF reification is not quotation
- The relata of the reified form are not names
- Reified form says what an RDF triple talks about, not what it is

http://www.w3.org/TR/rdf-mt/
RDF reification semantics - 2

- Intentionally weak entailment semantics ...
  - $T_{R \subseteq G_1} \not\equiv_{RDF} \exists(G_2)T \in G_2$
  - $T \in G_1 \not\equiv_{RDF} \exists(G_2)T_{R \subseteq G_2}$
- ... lead to identity problems
- To which $T$ in which $G$ does $T_R$ refer?
- Nothing in the language spec
- Requires semantic extension(s) to address

http://www.w3.org/TR/rdf-mt/
Operational issues

- SPARQL can query $T$ and $T_R$ in a single query
  - Assuming you’ve placed $T$ and $T_R$ in known graphs so they can be jointly accessed

- RDFS and OWL give you no way to do inference that involves reification

- You can refer to specific graph contents but not about their RDFS/OWL entailments
RDF reification – summary

• RDF reification lets you use sets of triples to talk about the subject matter of other triples

• You use extra-logical means (SPARQL) or semantic extensions implemented in your own code to connect the two

• RDF reification is blind to RDFS / OWL entailments
Reification in ECLIF
Reification in ECLIF

- Immediate goals same as RDF reification
  - Say things about asserted data
  - Query about those things
- Simple extension to ECLIF language
  - Syntax
  - Semantics
ECLIF reification syntax

Normal GAF

(likes bob pizza)

Reified GAF

(stmt s1 (likes bob pizza))

Fancier cases

(stmt s1 (between milwaukee chicago green-bay))

(stmt s1 (holdsIn (Date “2010”) (likes bob pizza)))
ECLIF reification semantics

- ECLIF has no “graphs” in its model theory
- More exactly, there is only one “graph”
- ECLIF reification thus differs from RDF
- For any ground atomic formula $F$
  - $F_R \models_{ECLIF} F$, or...
  - The reified form entails the non-reified GAF
  - Converse does not hold $\Rightarrow$ regress
When we assert a sentence, we mean to say that it is true.

We all know asserting it twice doesn’t make it “more true”.

This makes sense only if you’re only talking about uses versus mentions of a sentence.

Two mentions of a sentence could very well differ in their properties.

We call those separate mentions sentence tokens.
More types and tokens

• RDF includes the type/token distinction

• RDF reifications are intended to be references to **triple tokens**

• A token in RDF is the appearance of a **triple of a given form** in a graph

• This form is the **triple type**

• It is distinguished by the values of rdf:subject, rdf:predicate, and rdf:object
So, ECLIF reification is...

- ECLIF reification is a mechanism for naming tokens of ECLIF sentence types
  - Sentence types are not propositions
- It is possible to have multiple tokens of the same type in a single KB
  - Bags of sentence tokens (wrt type) vs sets
  - RDF semantics precludes this possibility
- Multiple applications
Life is simpler this way

- You’re probably reifying because you want to say something about your data
- You can simulate distinct RDF-like “graphs” by adding a graph-membership predicate
- Those simulated graphs will not segregate reified and object-level statements
- But they will allow you to keep statements in separate “containers” for applications
Simple example - stmt

Data

(stmt s1 (likes bob pizza))
(stmt s2 (likes bob beer))
(source s2 sue)

Queries

(stmt ?s (likes bob pizza)) => {{?s:s1}}

(stmt ?s (likes bob ?x)) => {{?s:s1,?x:pizza},{?s:s2,?x:beer}}

(and (stmt ?s (likes bob ?x))) => {{?s:s2,?x:beer,?source:sue}}
(source ?s ?source))
What’s new?

- Sentential operator – `stmt`
  - Used in ground data and in queries wherever atomic formulas may appear
  - Uses may not be nested
- Vocabulary for quantifying over statements
  - `Statement` - used to type and quantify over sentence tokens
  - `statementType` - used to map sentence tokens to sentence types
Summary

✔ Minimal ontologist effort to use
✔ Low storage, optimization & processing overhead
✔ Combined meta and object-level query
✔ No performance impact if reified data not used
✘ No HO quantification in current prototype
✘ Like RDF this is not real quotation
✘ No entailments
Proofs in ECLIF
Reification and inference working together
Problem statement

• Reification
  • Gives us names for sentence tokens that we can use to express metadata

• Inference
  • Generate “new” statements from asserted statements

• Reification + inference?
  • What would this look like?
Extra-logical approach

• Instrument query engine with machinery to construct proof objects alongside answer generation

• No necessary connection between
  • proof objects and domain entities
  • proof structure and ontology structure

• Return proof objects alongside answers

• This approach used by current XKS (≤ 4.x)
Logical approach

- Extend logic with meta-logical operators that construct proofs along with answers
- Proofs (along with statements) become part of extended UoD
  - Extension of reification approach
  - Can quantify over them → very powerful
- Prototype under development in XKS
Digression

• First we have to look at ECLIF’s properties as a logical system ...
• ECLIF semantics are WFS
  • Under non-mon Lloyd-Topor transform
• WFS is a logic programming semantics for interpretation of axioms in clause form
  • $H \leftarrow G_1 \land G_2 \land \ldots \land G_n$ ($n \geq 0$)
• Proofs will reflect this underlying logic
Syntax

Extend \texttt{stmt} operator to all sentences - not just GAFs

\[(\texttt{stmt} \ a1 \ (\rightarrow \ (\texttt{hasBrother} \ ?x \ ?y) \ (\texttt{hasSibling} \ ?x \ ?y)))\]

Sentential operator – \texttt{prove}

\[(\texttt{prove} \ ?p \ (\texttt{hasAncestor} \ p235 \ ?x))\]

May enclose any formula

\[(\texttt{prove} \ ?p \ (\texttt{and} \ (\texttt{hasParent} \ p235 \ ?x) \ (\texttt{hasBrother} \ ?x \ ?y)))\]

May be nested

\[(\texttt{prove} \ ?p1 \ (\texttt{and} \ (\texttt{hasParent} \ p235 \ ?x) \ (\texttt{prove} \ ?p2 \ (\texttt{hasBrother} \ ?x \ ?y))))\]
Proofs

• Proofs are trees
  • Non-leaf nodes are (ground) theorems
    • Unification (var instantiation) omitted
  • Leaf nodes are axioms (trivial theorems)
  • Non-leaf nodes result from application of proof operators to other nodes

• Each different proof represents a unique way of deriving a given answer - Key!
Example

simple ontology

a1: (<=
    (hasAncestor ?x ?y)
    (hasParent ?x ?y))

a2: (<=
    (hasAncestor ?x ?y)
    (and
      (hasParent ?x ?z)
      (hasAncestor ?z ?y)))

some reified data

(stmt s120 (hasParent p160 p70))
(stmt s137 (hasParent p189 p160))
(stmt s190 (hasParent p235 p189))
A sample proof

(prove ?p (hasAncestor p235 p70))

?p : (MP a2 (AND s190 (MP a2 (AND s137 (MP a1 s120))))))
How this works

• **Class of proofs** – **Proof**
• **Reified statements prove themselves**
  • **Statement** is a subclass of **Proof**
• **Proof operators construct proofs**
• **Proofs can be quantified over**
  • Their parts can be inspected
  • Parts that are statements carry metadata
Proofs and metadata - recipe

(and (prove ?p (hasAncestor p235 p70))

formula for proof analysis)
Application examples

Multi-level security filtering

(and (prove ?p (hasAncestor p235 p70))
  (proofSecLevelLTE ?p SecLevel_S))

Provenance

(and (prove ?p (hasAncestor p235 p70))
  (exists (?a)
    (and (axiom ?p ?a)
      (hasSource ?a BaltimoreSun))))

Probability / confidence

(and (prove ?p (hasAncestor p235 p70))
  (proofConfidence ?p ?conf))

Proof complexity

(and (prove ?p (hasAncestor p235 p70))
  (count (and (subProof ?p ?s) (= ?s (MP ? ?)))
    ?count))
Summary

✔ Metadata with inference
✔ No ontology impact
✔ Processing overhead ≈ proof complexity
✔ Combined meta and object-level queries
✔ No performance impact if proofs not needed
✘ Proof identity an issue – order dependence

✘ Requires commutative algebraic operators
Questions?